

# An Innovative SDR Architecture and MANET with Simultaneous Multiple Channel Reception

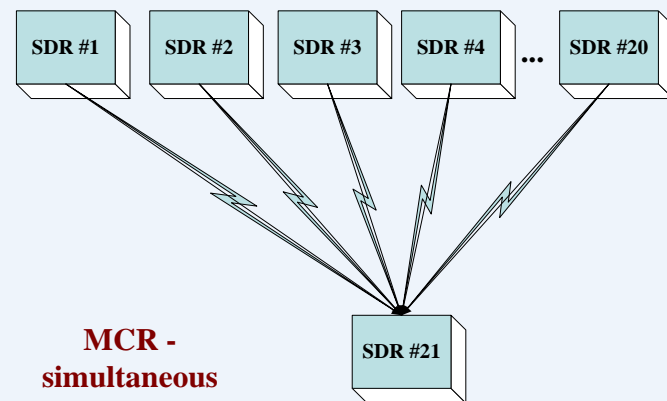
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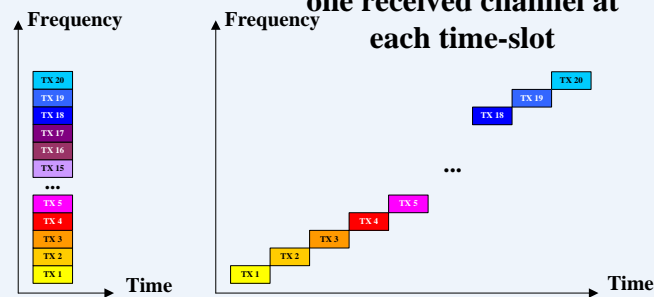
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**MCR - simultaneous reception of > 20 independent channels**

**Other modern MANET – one received channel at each time-slot**



## Presentation Summary

- **Tactical MANETs usually employ 1(2) Hw receivers (per band), thus they are challenged by –**
    - ✓ **Tactical communication is mostly asymmetric, Rx rate  $\gg$  Tx rates - Receiver Bottleneck**
    - ✓ **Besides “Intra-Unit” nets, “Inter-Unit” geographical Network Scalability is needed (for SA, FA, CA)**
    - ✓ **Units frequently Merge, pass-thru, Split - need to avoid reconfiguration.**
- ➔ *Dynamic, High Scalability needed***
- **The Solution: Multiple-Numerous-Channel Reception architecture, based on Ultra WB ADC, Ultra Fast FPGA**
    - ✓ **Hundreds channels received (per band)**
    - ✓ **Almost unlimited Scalability, > 200 Mb/sec Rx rate!**

# Scalability Challenge of 1-Channel TDMA Networks – Merging, Passing-thru and Splitting forces

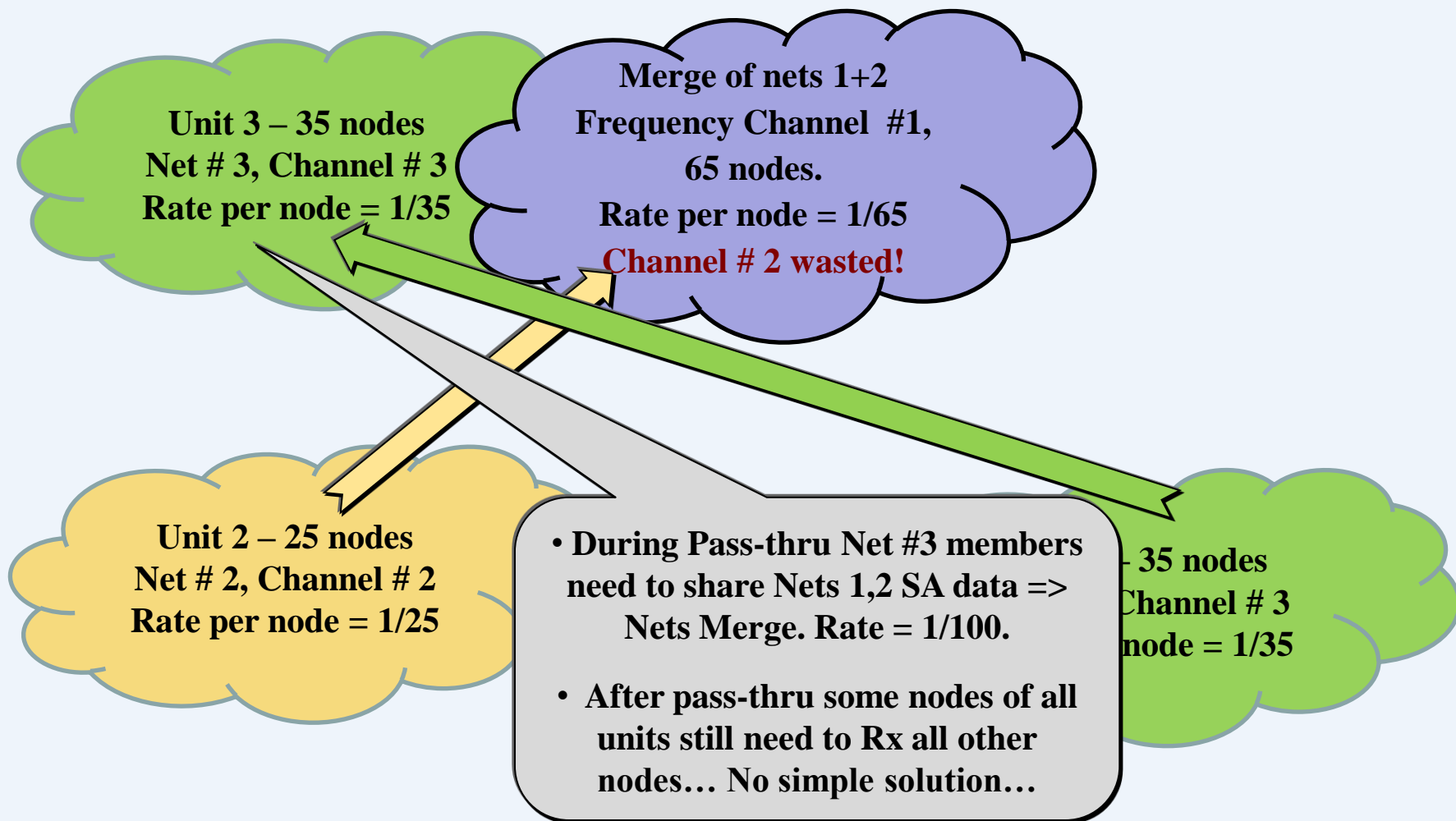
**Unit 1 – 40 nodes**  
**Net # 1, Channel # 1**  
**Rate per node = 1/40**

**Initially, units are Isolated => Cross (Inter)-unit update rate can be low**

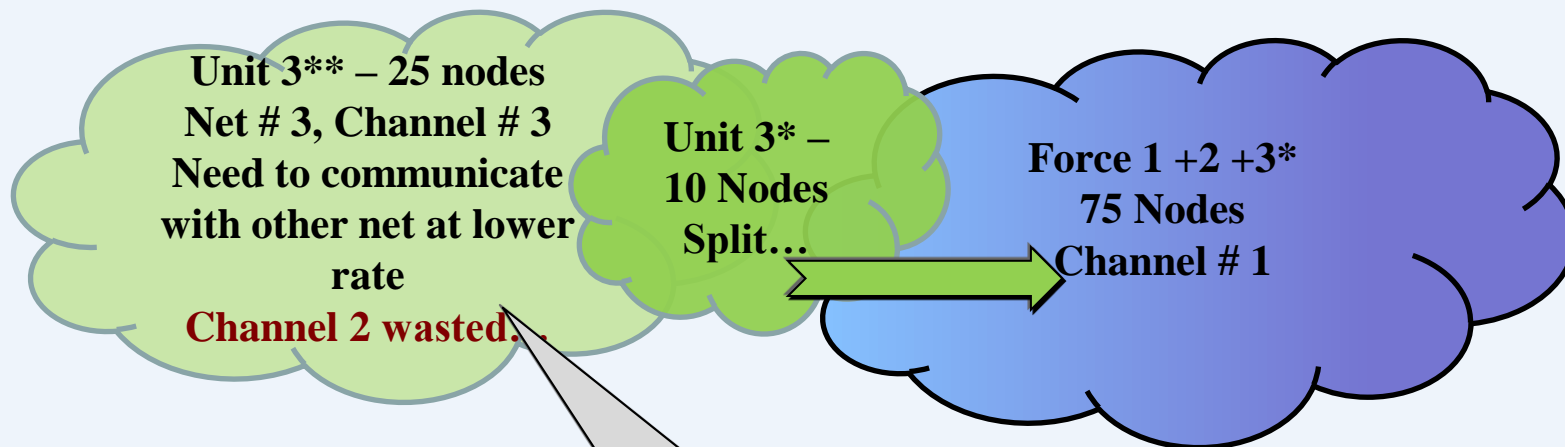
**Unit 2 – 25 nodes**  
**Net # 2, Channel # 2**  
**Rate per node = 1/25**

**Unit 3 – 35 nodes**  
**Net # 3, Channel # 3**  
**Rate per node = 1/35**

# Scalability Challenge of 1-Channel TDMA Networks – Merging, Passing-thru and Splitting forces (2)



# Scalability Challenge of 1-Channel TDMA Networks – Merging, Passing-thru and Splitting forces (3)



• After pass-thru, 10 nodes from U3 merge with 1+2. Some nodes of all units still need to Rx all other nodes... No simple solution...

# Scalability Challenge of 1-Channel TDMA Networks – Merging, Passing-thru and Splitting forces (4)

Unit 3\*\* – 25 nodes

- ✓ Allocation of channels to Intra-unit networks demands frequent reconfiguration in dynamical combat scenarios.
- ✓ “Merge/Split” algorithms – complicated, slow response, consume control bandwidth....
- ✓ SA/CA/FA Throughput of merged network don't scale - decreases linearly with N....
- ✓ Some available frequency channels can't be used....

**Simultaneous, Multiple Channel Reception solve all these problems!**



# Scalability Challenge of 1-Channel TDMA Networks: “Receiver Bottleneck” Effect



Total Rx rate 1 Mb/sec

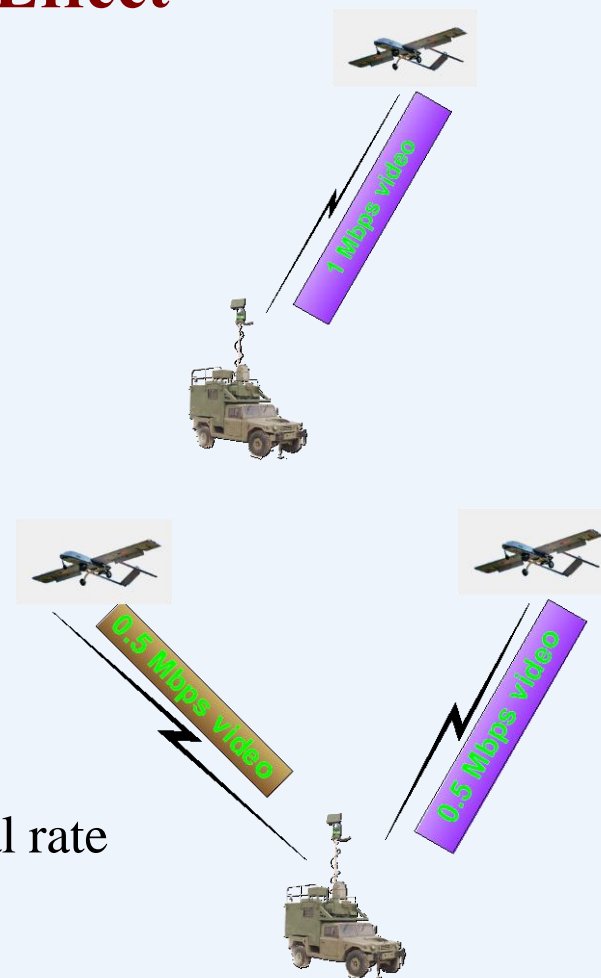


Total Rx rate is still 1 Mb/sec..

Reception of 2 Video Sources: Time Sharing.

## Tradeoff:

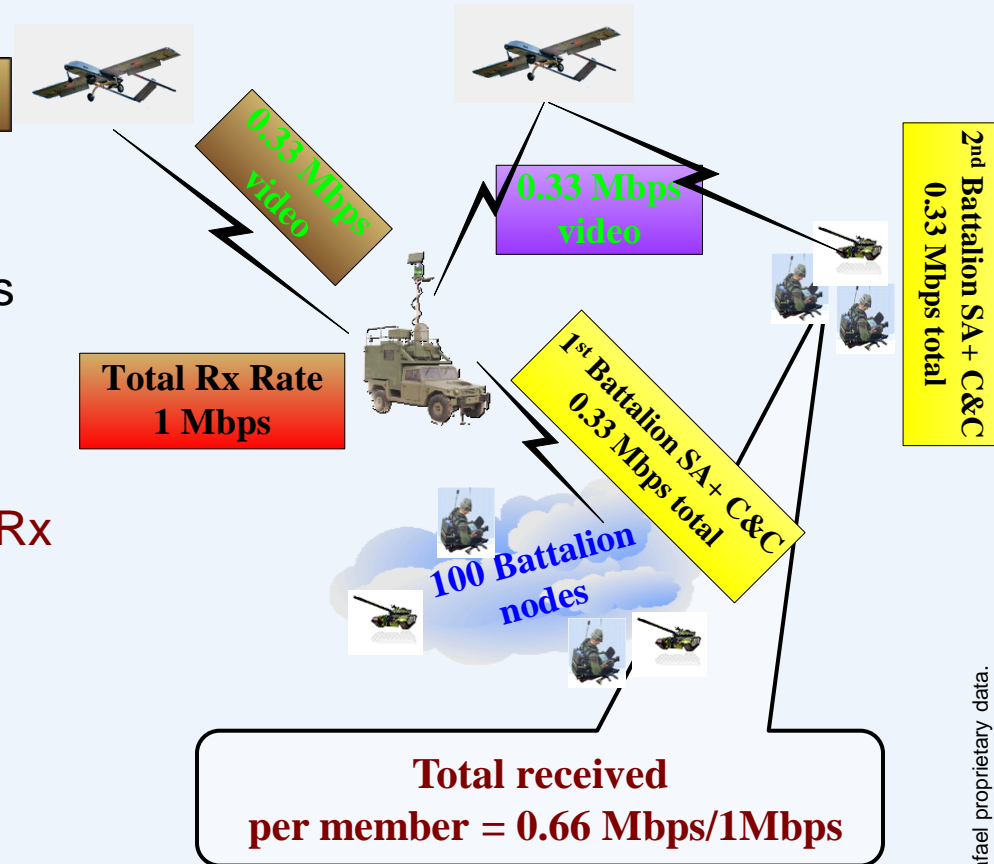
- Each source should be transmitted at 50% of total rate



# Scalability Challenge of 1-Channel TDMA Networks: “Receiver Bottleneck” Effect (2)



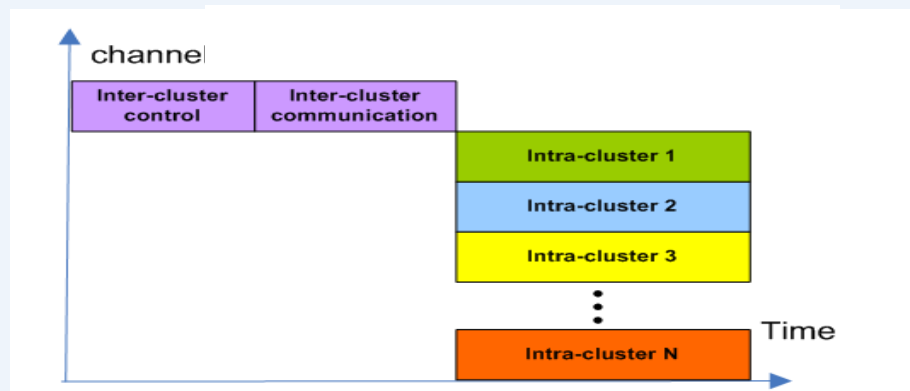
- 1<sup>st</sup> HQ Rx 2 Video Sources + 100 nodes Battalion SA + C&C messages
- 2<sup>nd</sup> Battalion need only 1 Video + Battalion SA + C&C – Could Rx 0.66 Mbps Video, but is limited by 1<sup>st</sup> HQ Rx bottleneck...
- **Even More Tradeoffs...**



**Throughput is bounded by the Bottlenecked Receiver, not by Tx...**

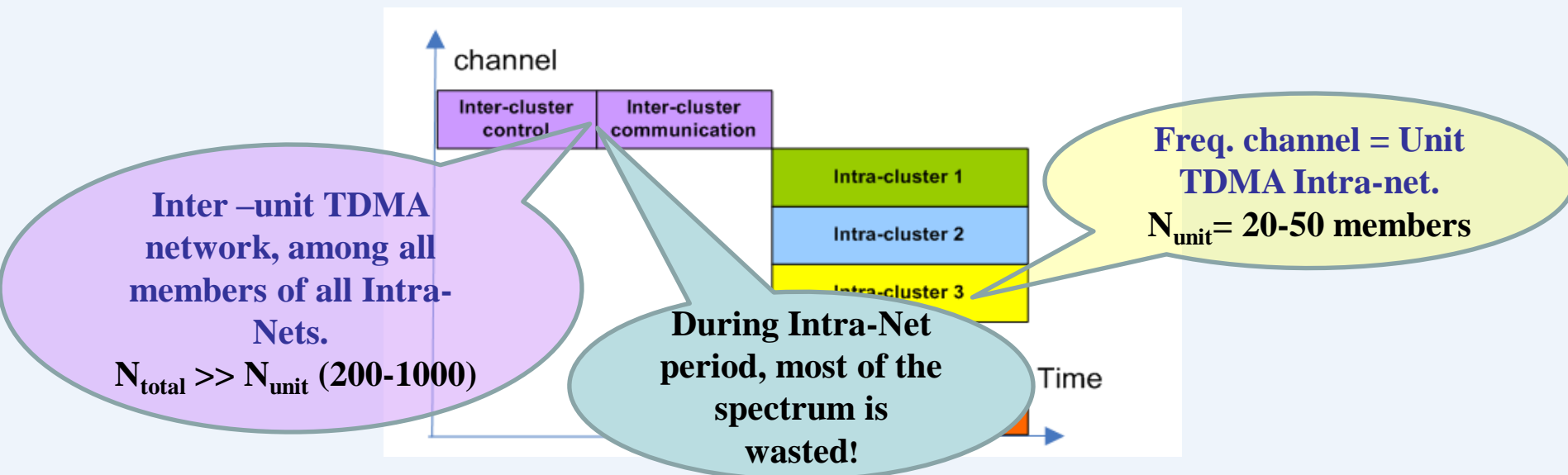


## A Common non-scalable approach - “Multiple 1-Channel TDMA” MANET



- ✓ **Most SDRs can receive only one channel per band at a time.** Once frequency has been set, the networking is done only among nodes tuned to that channel. The MAC is “blind” to all other channels.
- ✓ **A MANET system with a collection of channels, is in fact a collection of unconnected parallel MANETs,** with data rate limited by the BW of a single channel.
- ✓ **Time- Sharing between Intra/Inter- Unit Communication. Inter-Unit is done over a single channel. → very low rate and waste of most of the spectrum during the Inter-Network time frame.**

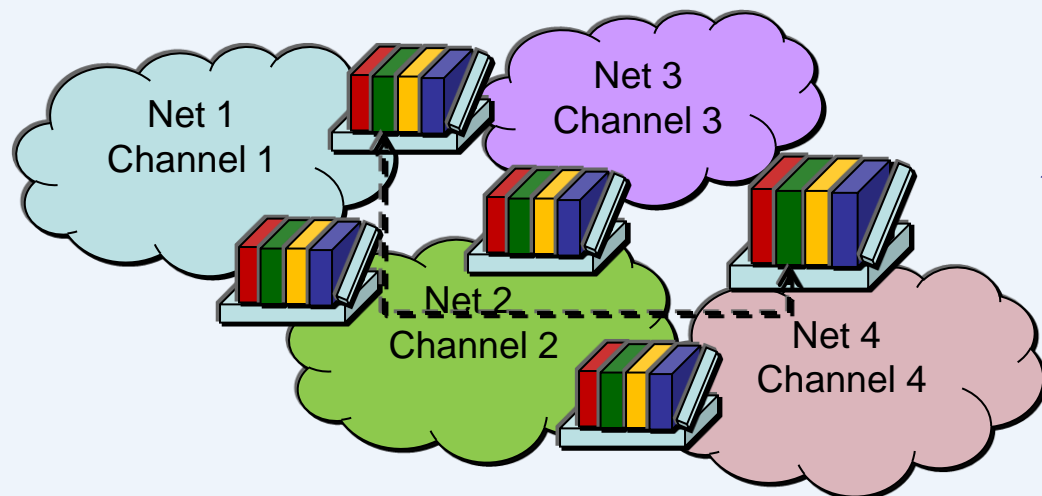
## “Multiple 1-Channel TDMA” approach



- Each radio can receive only a single channel (per band) simultaneously
- “Freq. Channel = Intra-unit TDMA net” - A/C formation, Company, Battalion...
- “Inter-Unit” communication is done over a Bottlenecked time shared single channel
  - ✓ Waste of spectrum!
  - ✓ Very low update rate per each user in the Inter/ Cross-Unit network
  - ✓ ➔ Inadequate SA/CA rates for merging units/formations.

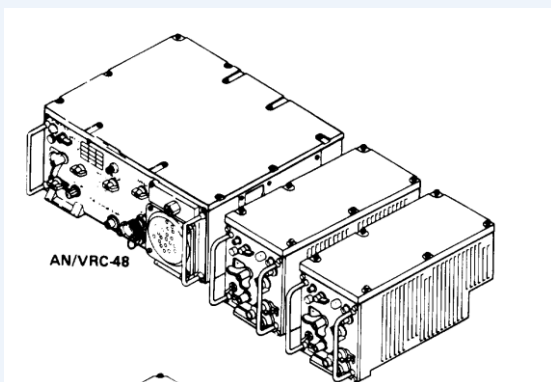
## Another Low Scalability Approach – Multi Transceiver Inter-Net Gateways

- ✓ Each sub-net has several multi- transceiver SDRs which serve as gateways
- ✓  $A(f1) \Rightarrow \text{Gateway}(f1) \Rightarrow (\text{thru Lan}) \text{Gateway}(f2) \Rightarrow B(f2)$



- ✓ Lots of expensive, heavy radios.
- ✓ If a Sub-net has  $L > 4$  neighbor Sub-nets, it needs  $L$  hardware transceivers...
- ✓ **Throughput still very low!** – every packet has to be retransmitted in all destination networks.

## Historical Reminder....



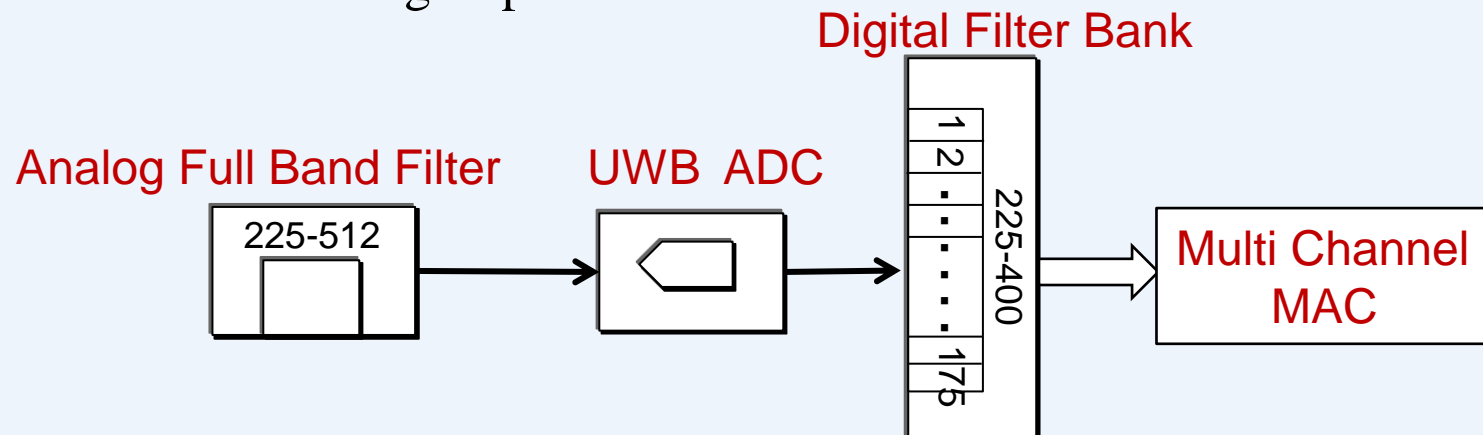
### VRC-44: The Commander's Radio (1965-1990+..)

- ✓ 1 Transceiver + 2 “Auxiliary receivers”.
- ✓ Can Rx 3 unit or “functional” networks.
- ✓ Transceiver used for “Network Entrance”.
- ✓ Command Vehicle/Post had 3-5 sets – Commander, Int., Fire, Logistics....

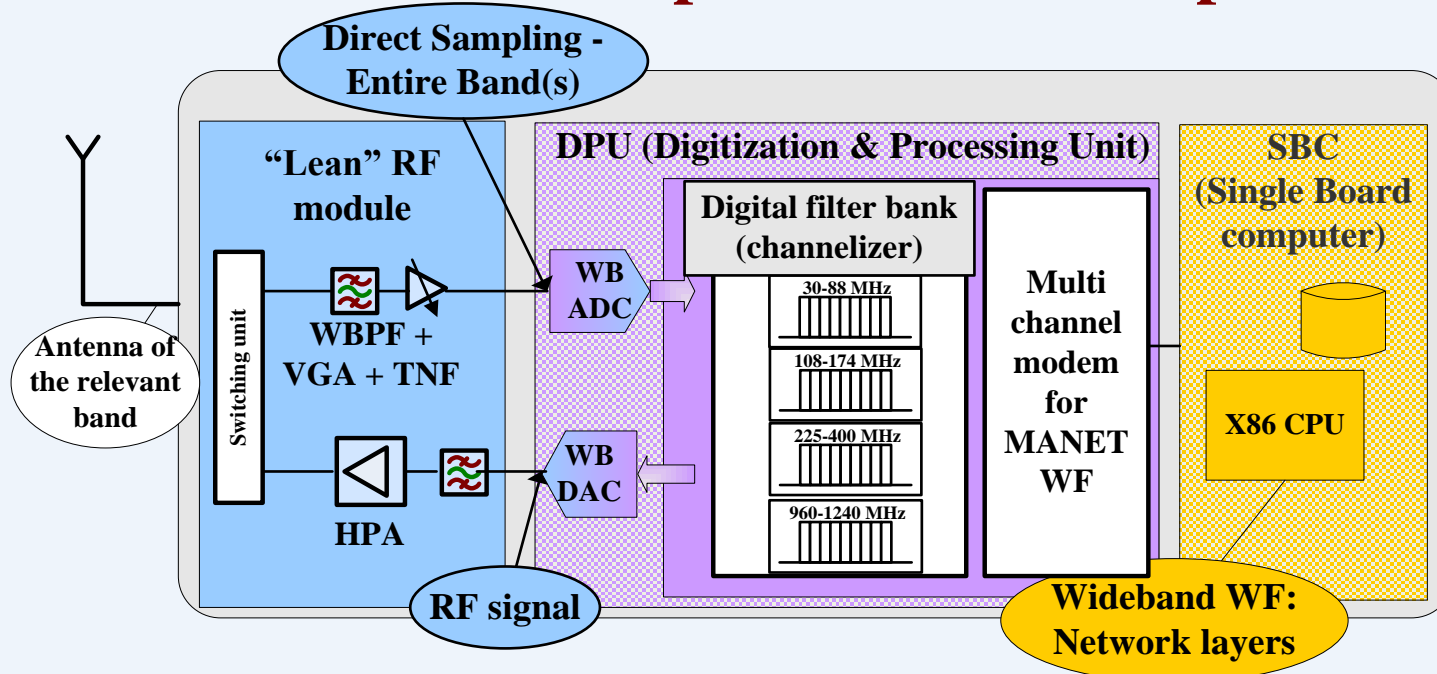
- **Ideally, A Military MANET SDR should have numerous receivers** - “Listen” to all desired Unit and Functional networks in it’s neighborhood, Tx when it needs.

## Multiple Channel Reception MANET SDR

- Preferably, A Military MANET SDR should have numerous receivers- “Listen” to all desired Unit and Functional networks in it’s neighborhood, Tx when it needs.
- **Multi-Channel digital Reception (MCR) architecture -**
  - ✓ Direct sampling of a whole band(s) - Almost an “Ideal True SDR”, **achieving 10÷100 fold increase in Rx throughput..**
  - ✓ Tx frequency of a node is arbitrary in the allocated band, not related to unit or functional group.



# Ultra Wide Band Multiple-Channel Reception SDR



- ✓ **Ultra Fast/ UWB ADC** – Dual 1.8 GSPS/12 bits or 1.2 GSPS/14 bits.
- ✓ **Ultra Fast FPGA** – more than 1 Tera-Ops/sec.
  - **Network ≠ Frequency channel.** Networks are *logical*. Channels allocated auto. by MAC.
  - **No reconfiguration during merge/pass thru – all nodes listen to all channels**
  - Tx node doesn't coordinate frequency with Rx nodes. Freq. Hoping – in Tx only...
  - **“Cognitive Ready”** – Rx is an instantaneous UWB Spectrum Analyzer...



## Tactical MANET Periodical SA Messages (PM)

- **Basic service for C&C applications** (e.g., “Blue Force Tracking”)
- **Periodical (refreshing) short messages – Multicast/Bcst packets for-**
  - ✓ **Situation Awareness (SA)** – Platform/Unit position/speed, status, tracked enemy targets– for COP and FA (Fratricide Avoidance).
  - ✓ **Collision Avoidance (CA)** – very high rate (10 Hz) needed in Airborne platforms; position and speed.
  - ✓ **Control packets** – rapid dissemination of network topology.
- **Destination – All nodes in relevant geographical area, regardless of unit association**
- **High Scalability needed while keeping high update rate**, in dynamic combat scenarios, Companies, Battalions, A/C formation – frequently merge or pass-thru.
- Need a scalable, fast and simple merging process w/o network reconfiguration.

# PM-MCR MAC: providing a true scalable solution for broadcast PM data

## ○ The Ideal (static F/TDMA) Case -

- ✓ **N nodes** fully connected. **M frequency channels**. All nodes know IDs of all others.

**CHANNELS**

1	46	..	..	24	62	..	..	8	16	14
27	13	..	..	58	32	..	99	37	42	49
78	5	..	..	83	1	..	56	99	71	..
14	100	..	..	8	91	..	..	27	78	..
99	66	..	..	..	70	..	..	51	5	..

← **N/M Time Slots** →

Nodes 1, 27, 78, 14, 99  
Don't receive each other on 1<sup>st</sup> time slot. Random permutation of Tx schedule will enable receiving each other in following blocks.

- ✓ **Half-Duplex Radio**, hence all radios Txing on same time-slot don't Rx each other. ⇒ a small rate loss of  $(M-1)/(N-1)$  – small for  $M \ll N$ .
- ✓ Tx order of every Block of  $N/M$  time-slots is a random permutation of the node IDs, with Seed = f(Block #).
- ✓ Synchronized network, all nodes execute same permutations with identical seeds.

## Static PM-MCR (F/TDMA) Performance

- The average Broadcast messages reception rate (every node Rx any other node) –

$$(1) \quad R_{PM-MCR} = \frac{M}{NT_{slot}} \left[ 1 - \frac{M-1}{N-1} \right] = M \left[ 1 - \frac{M-1}{N-1} \right] R_{1Channel TDMA}$$

$$\xrightarrow[M \ll N]{} MR_{1Channel TDMA}$$

“Half-duplex  
Loss”

- **Average Throughput Gain of a MCR SDR w.r.t. 1-channel TDMA SDR, is  $\sim M = \#$  of Channels....**

- ✓ Compared to “M- Multiple/Stacked 1-TDMA” Inter-net, where  $X < 1$  of the time-frame is allocated to the Inter-units network, ATG is  $1/(1-X)$  higher than (1).
- ✓ **Example:**  $N = 501$  nodes,  $M = 21$  Channels,  $X = 0.5$  of time frame for Inter-Unit net  $\Rightarrow$  **ATG** of the Inter-Unit communication w.r.t. “M-stacked 1-TDMA” = **40.3**

## Static PM-MCR (F/TDMA) - continued

- Due to random permutations of node's  $k$  "transmission cell" within a block of  $\{M \text{ frequency channels} \times (M/N) \text{ time-slots}\}$ , and the random # times node's  $k$  transmissions are not heard by some other nodes transmitting at the same time slot with  $k$ , the probability of reception is Binomial –

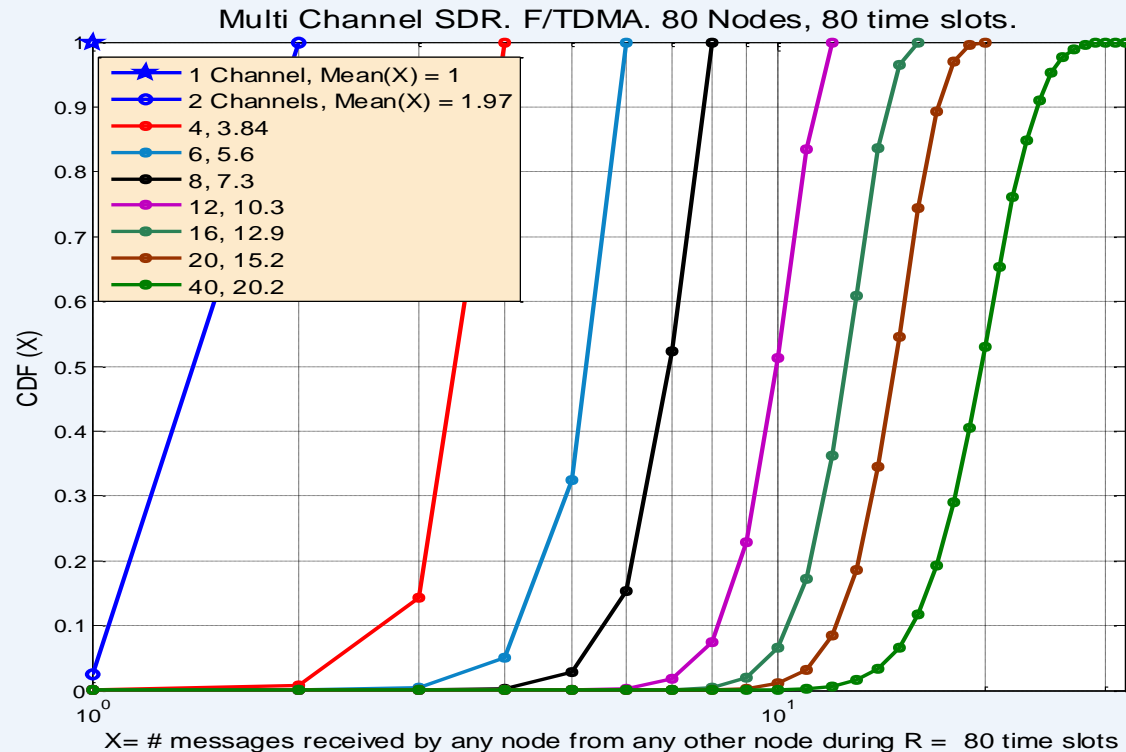
$$(2) \quad \text{Pr}x_{ijL}(k) \triangleq \text{Prob}\{\text{Node } i \text{ received node } j, k \text{ times during } R \text{ time - slots}\}$$

$$= \binom{L}{k} \text{Pr}x_1^k (1 - \text{Pr}x_1)^{L-k}; \quad k = 0, 1, 2, \dots, L; \quad L = \frac{RM}{N}$$

$$(3) \quad \text{Pr}x_1 = 1 - \frac{M-1}{N-1}$$

- The CDF however rises sharply  $\Rightarrow$  "Fairness".
- Simulations show that the theoretical CDF is achieved within 2-3 frames of  $(N/M)$  Time-slots.

## Static PM-MCR (F/TDMA) - continued



- **Discussion:** For **80 Nodes, 200 Time-slots/sec**, if **4 channels** are allocated for PM service, **Average reception rate =  $2.5 \times 3.8 = 9.6$  Messages/sec**, compared to **2.5 Mess/sec with a 1-Channel TDMA**.
- More than  $2.5 \times 2 = 5$  Mess/sec are received with 100% probability.
- Only 10% are received at rate lower than 8.5. Mess/sec

## PM-MCR almost achieves Shannon's bound for 1-Channel UWB TDMA (Full Mesh case)

- **Objective:** Show optimality of MCR F/TDMA – it almost achieves Shannon's Bound for a 1- channel UWB TDMA, with the same BW of M- Channels MCR.
- **However, practically MCR F/TDMA MANET is always better than a 1-channel UWB TDMA, since UWB requires –**
  - ✓ Very high peak power ( $P_{\text{uwb}} = M * P_1$ )
  - ✓ A wide contiguous spectrum allocation ( $B_{\text{uwb}} = M * B_1$ )
  - ✓ Packet length  $T_{\text{UWB}} = T_1 / M$  is too short for allowing Guard Time for long range, airborne communication.
    - (e.g.: 100 Km range  $\Leftrightarrow$  0.33 Millisec GT. Packet size = 1 Kbyte  $\Leftrightarrow$  1 Millisec at 1 Mb/sec, 0.025 Millisec at 40 Mb/sec  $\ll$  0.33.....)
- **Proof Outline –**
  - ✓ Requiring Average Tx power is equal for both cases  $\Leftrightarrow$  Equal Range
  - ✓ Requiring # of bits/packets (neglecting Guard Time) is equal in both cases



# PM-MCR almost achieves Shannon's bound for 1-Channel UWB TDMA

## Definitions:

**N**- # Nodes. **M**- # channels.  $B_u = M \cdot B_1$  - BW of UWB radio.  $T_1, T_u$  - Time slot length. **P**- peak

### Multi Channel SDR

### UWB 1-Channel SDR

#### (4) Shannon's Bound

$$C_1 = B_1 \text{Log}(1 + SNR_1) \text{ per channel.} \quad C_{uwb} = MB_1 \text{Log}\left(1 + \frac{P_u}{MP_1} SNR_1\right) = MC_1$$

Requiring Average Tx power is equal for both  $\Rightarrow P_{uwb} = M \cdot P_1$

(This condition also implies identical SNR, thus identical reception range)

(5) Node's Message Rx Rate, from eq. (1) -

$$R_{PM-MCR} = \frac{M}{NT_1} \left[ 1 - \frac{M-1}{N-1} \right] \quad R_{uwb} = 1/(NT_u)$$

Requiring # of bits/packets,  $C \cdot TS$  (neglecting Guard Time), is equal in both cases:

(6)

$$\frac{T_1}{T_u} = \frac{C_{uwb}}{C_1} = M$$

## PM-MC almost achieves Shannon's bound for 1-Channel UWB TDMA (2)

Substituting (6), (5) into (4), the message rate ratio is –

$$\frac{R_{PM-MC}}{R_{UWB}} = \frac{MNT_u}{NT_1} \left[ 1 - \frac{M-1}{N-1} \right] = \left[ 1 - \frac{M-1}{N-1} \right] \xrightarrow{M \ll N} 1$$

- ✓ This theoretical result shows that PM-MCR (for static, fully connected mesh case) almost achieves Shannon's bound.
- ✓ However, practically MCR MANET is always better than a 1-channel UWB TDMA, since UWB requires –
  - ✓ Very high peak power ( $P_{uwb} = M \cdot P_1$ )
  - ✓ A wide contiguous spectrum allocation ( $B_{uwb} = M \cdot B_1$ )
  - ✓ Packet length  $T_u = T_1/M$  is too short for allowing Guard Time for long range, airborne communication.  
(e.g.: 100 Km range  $\Leftrightarrow$  0.33 Millisec. Packet size = 1 Kbyte  $\Leftrightarrow$  1 Millisec at 1 Mb/sec, 0.025 Millisec at 40 Mb/sec  $\ll$  0.33.....)

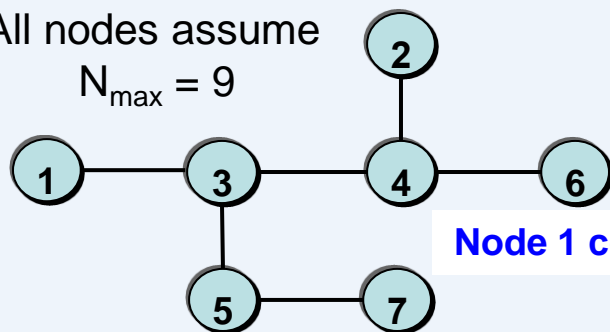
## PM-Multiple Channels Adaptive distributed F/TDMA MAC

- **Previous example was for static mesh-** not adequate for high mobility, nodes are not aware of all other active nodes in the whole Inter-network.
- **Adaptive PM MAC for MCR** – A 2 Dimensional generalization of a common class of “Adaptive TDMA” algorithms (e.g. MALS, SEEDX, NAMA alike –[1-4]), where -
  - ✓ There is a “Network joining” (and leaving) process, during which a new node learn about its 2-3 hops “Contention Neighborhood” (all nodes that could contend with it directly or interfere with its destinations as “Hidden Terminals”).
  - ✓ For each Time/Frequency cell (TFC) a distributed election algorithm is used to select one transmitter, according to a deterministic “Priority function”, which induces a deterministic order among all contending nodes in a neighborhood, by means of deterministic function (e.g. a permutation of {Nodes ID Vector}, PRNG, Latin Square..). This function is based on a “seed” =  $f(\text{TFC})$  known to all contenders.
  - ✓ This election ensures that all nodes in the one-hop neighborhood of the transmitter will receive data without any collision. The seed pseudo-randomly changes in each TFC, thus ensuring (some) fairness among nodes.

# Example of an Adaptive distributed TDMA Algorithm

Scheduling matrix – PRN permutations of {1,2,...9}

All nodes assume  
 $N_{max} = 9$



Node 1 contenders: {3,4,5}

TS 1	6	3	7	8	5	1	2	4	9
TS 2	7	2	8	5	6	9	4	1	3
TS 3	4	3	7	9	8	1	5	6	2
TS 4	5	7	8	3	6	1	2	4	9
TS 5	4	9	2	6	5	1	7	8	3
TS 6	9	6	2	1	3	8	7	4	5
TS 7	1	7	8	4	2	9	5	6	3
TS 8	5	3	9	7	4	1	8	2	6
TS 9	3	2	6	4	9	8	7	1	5
TS 10	9	2	1	6	8	3	7	4	5

**Fairness depends on –**

- Type of “Order Inducing Function”
- $N_{max}$
- “Fairness window length”

**Optimizations possible...**  
 e.g., LS is better, but more computationally demanding for large  $N_{max}$

# Tx in 10 TSlots	Node	Contenders	Time Slots									
			1	2	3	4	5	6	7	8	9	10
3	1	3,4,5	3	5	4	5	4	1	1	5	3	1
2	2	3,4,6	6	2	4	3	4	6	4	3	3	2
1	3	1,2,4,5,6,7	6	7	4	5	4	6	1	5	3	2
2	4	1,2,3,5,6	6	2	4	5	4	6	1	5	3	2
2	5	1,3,4,7	3	7	4	5	4	1	1	5	3	1
2	6	2,3,4	6	2	4	3	4	6	4	3	3	2
2	7	3,5	3	7	3	5	5	3	7	5	3	3
<b>Tx Winner</b>			6	2	4	5	4	1	1	5	3	1
				7				6	7			2

Notice that algorithm achieves some Spatial Reuse

# PM-Multiple Channels Adaptive F/TDMA MAC

## ○ Extension to Multi Channel SDR:

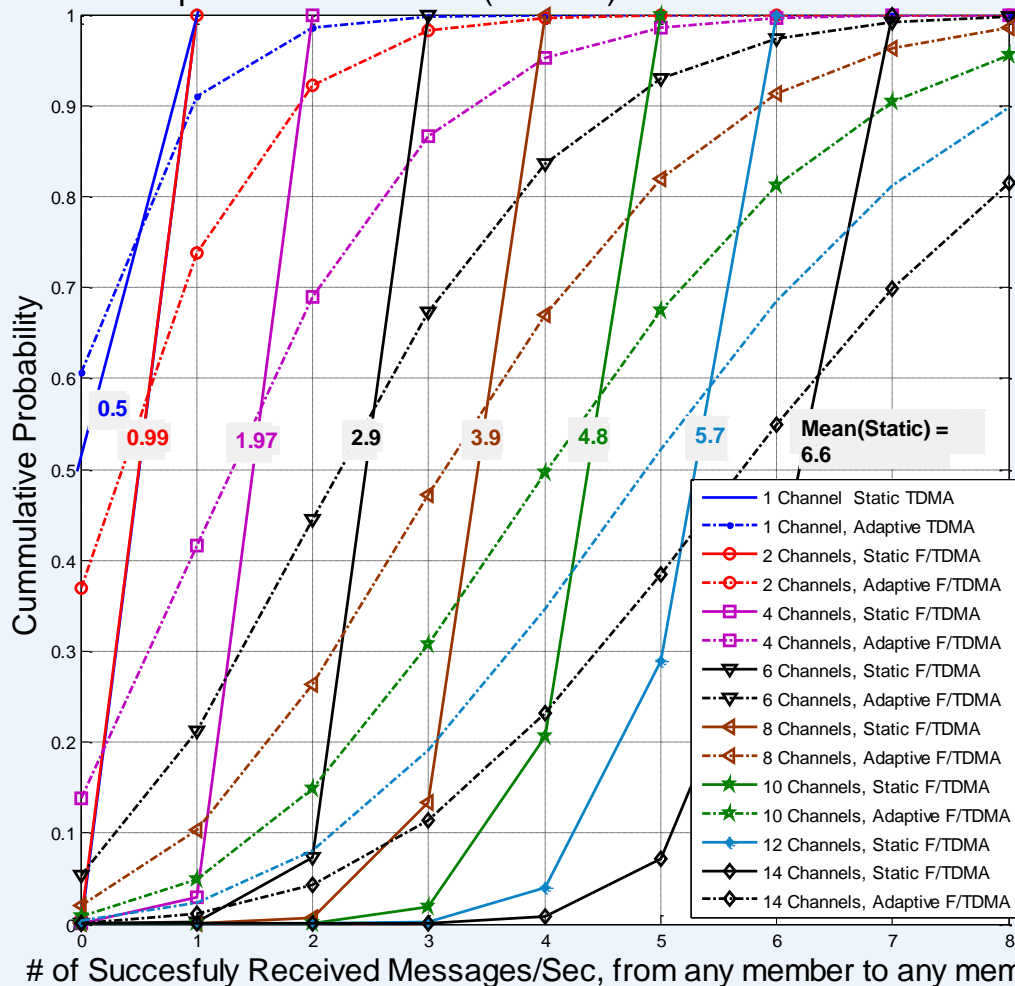
- ✓ The 1-channel Adaptive TDMA algorithm is run per each Time/Frequency Cell (TFC) with a PR seed =  $f(\text{TFC})$ , known to all nodes. The node with the highest priority among it's known contenders elect itself to transmit.
- ✓ If a node wins 2 TFCs in a Time-Slot, it transmits in only one of them (\*).

## ○ Optimization allows QOS and better fairness –

- ✓ Type of “Order Inducing Function” (Latin Square, PRN, Hash...)
- ✓ Total # IDs (Nmax) w.r.t. actual network size.
- ✓ When a node needs higher throughput – it is given several IDs.

# “Neighborhood aware” PM-MCR Adaptive F/TDMA MAC

SDR with MCR, CDF of Update Rate, 100 Slots/Sec, 200 Members  
Adaptive F/TDMA MAC (dashed) Vs Static



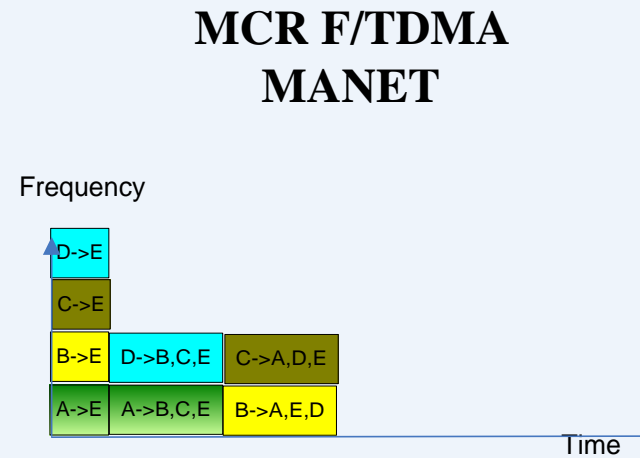
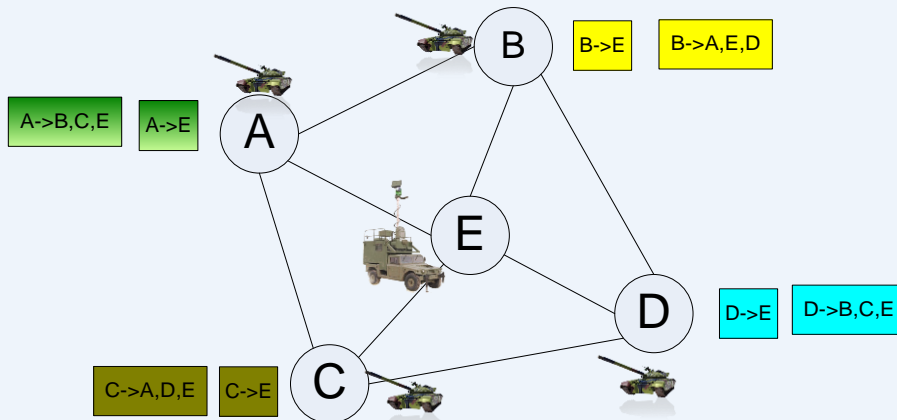
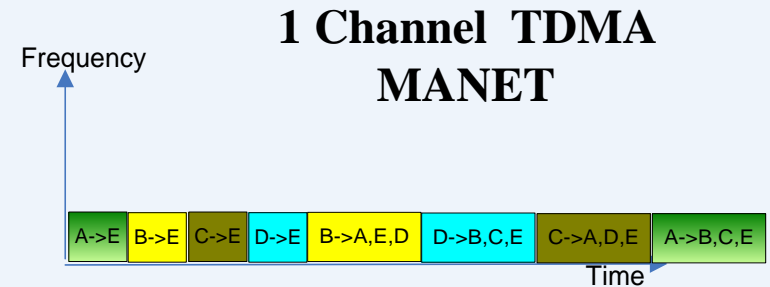
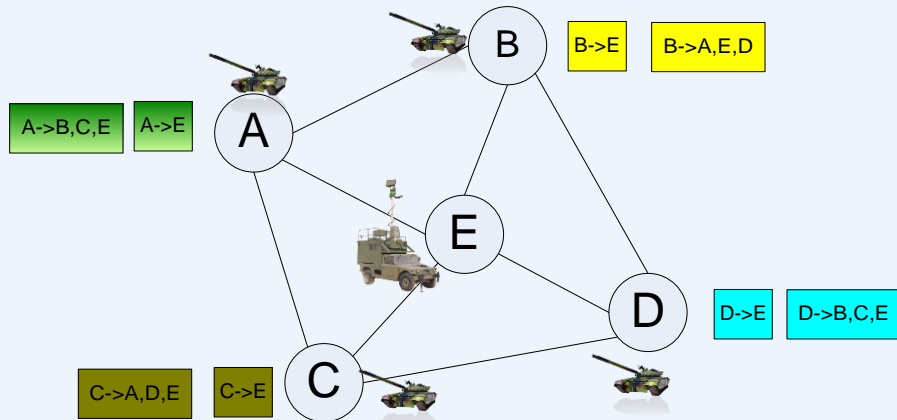
- ✓ Full lines – Static MAC
- ✓ Dashed lines – simulation of adaptive F/TDMA
- ✓ All nodes are fully connected
- ✓ # Channels to ensure 3 Mess/Sec:
  - “6.2” Mean (Eq. 1, roughly M times vs. 1 -Channel SDR)
  - **8 Ch. at 50%**
  - **10 Ch. at 70%**
  - **12 Ch. at 80%**
- ✓ “Full Mesh” is worst case – in other scenarios Throughput & Fairness are better due to Spatial Reuse

**Full- Mesh is worst case Topology (better spatial-reuse in other scenarios)**

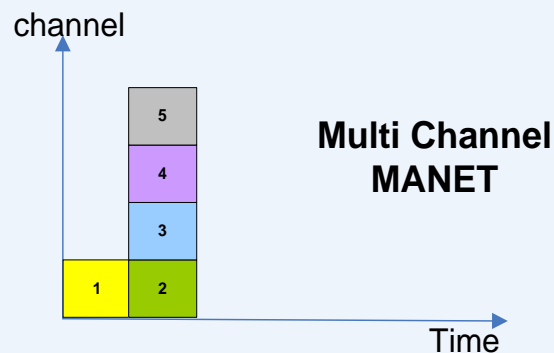
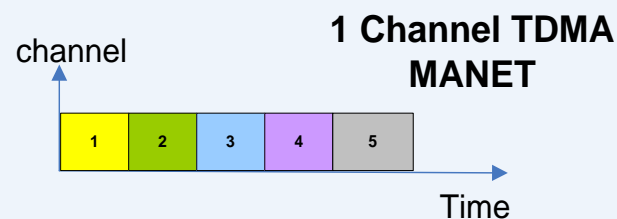
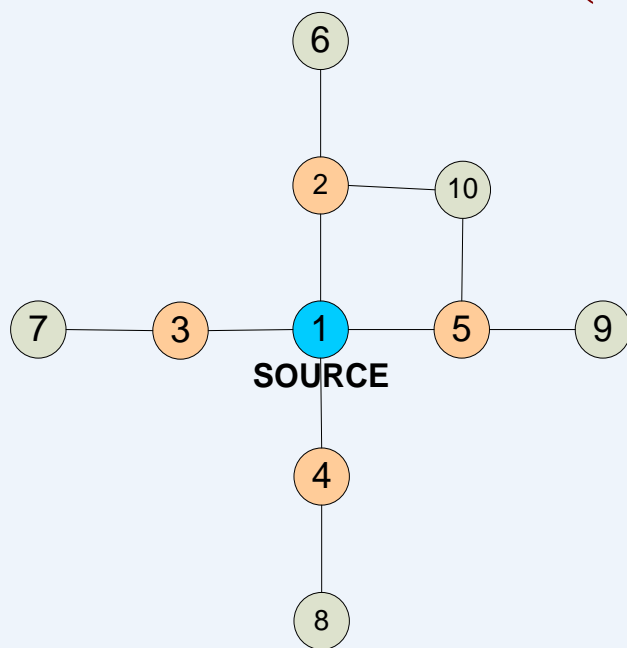


# Network Delay- arbitrary length Unicasts /Multicasts

## The bottleneck of 1 Channel Manets is the receiver....

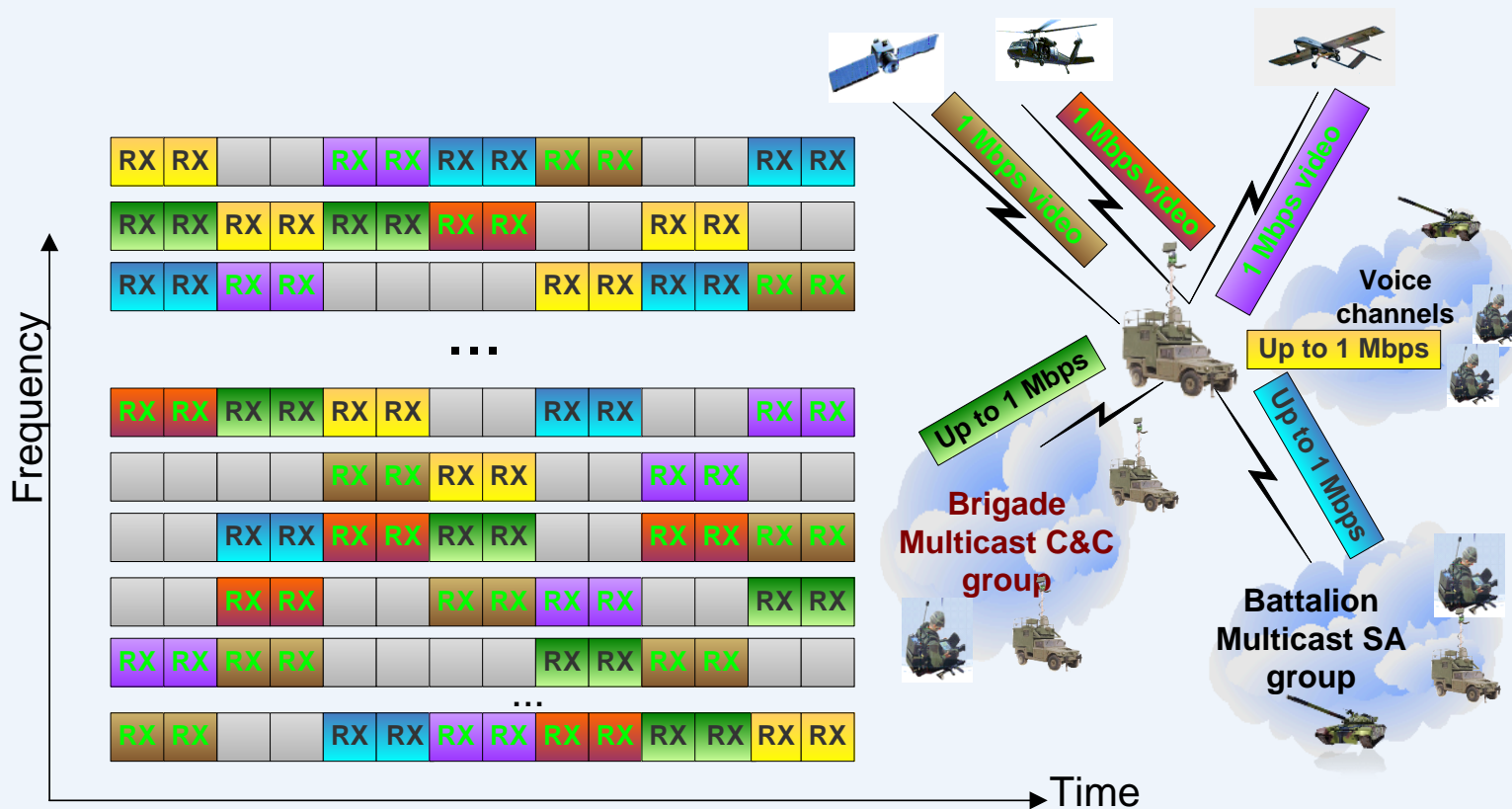


# Shorter Delay when Routing delay critical services (Voice, Video)



- **MCR - No delay**
- **1 Channel - 3 frames max, 1.5 frames average**

# MCR F/TDMA – Combining different Comm. Services

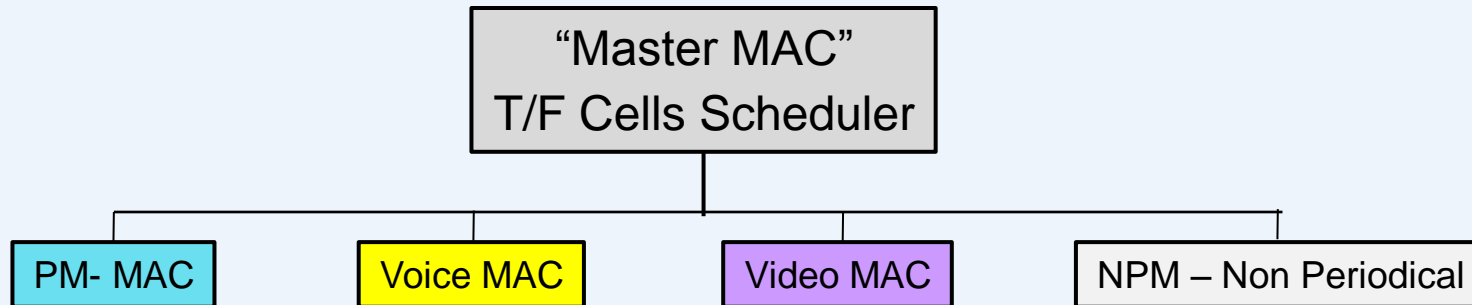


Reception of several frequency-hopping channels with **No Tradeoff !!**

**Total received data ~ 6 Mb/s ~ Sum of Tx rates ⇒ (# Channels)X Channel rate**

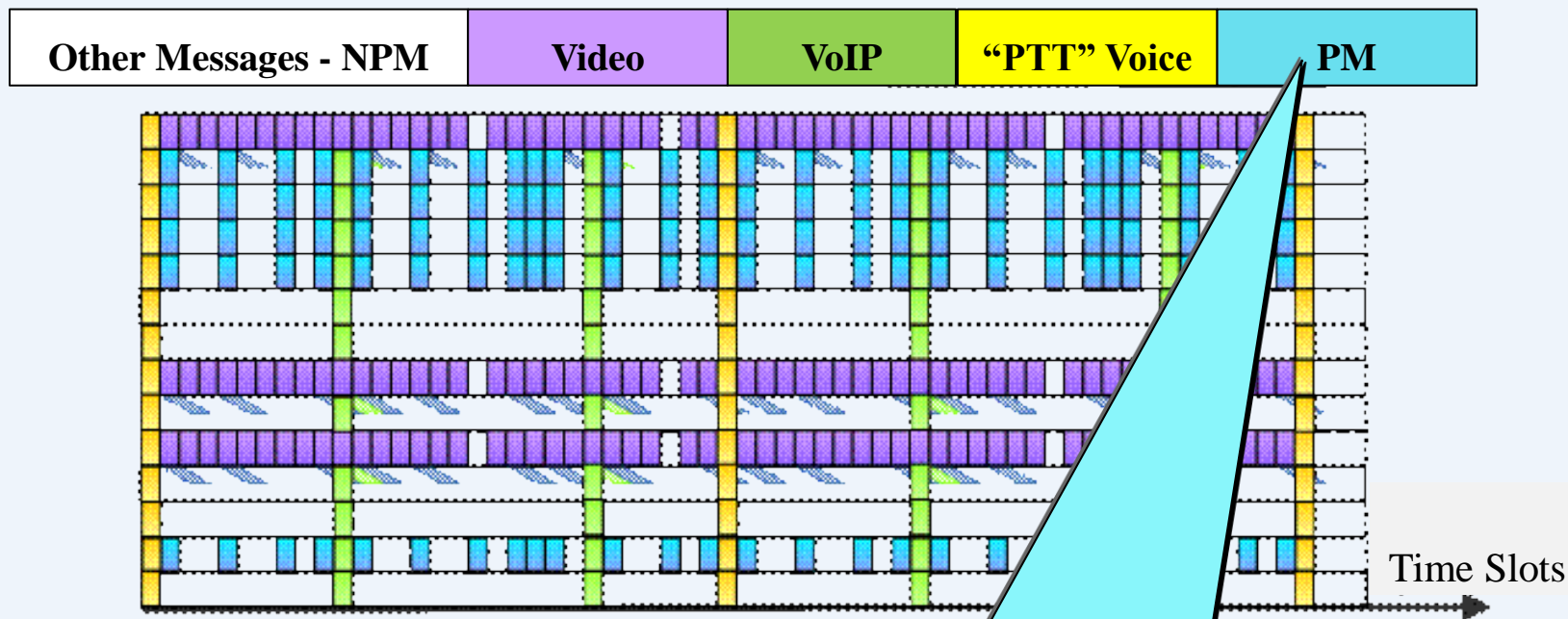
**The Data Rate is bounded only by the rate of the transmitters**

## MACs for different communication services



- **Each service type has different QOS and Routing/ Relaying requirements -**
  - ✓ Broadcast, Multicast, Unicast.
  - ✓ Delay critical or latency tolerant.
  - ✓ Broadcast Relaying (for “Combat Voice”) or N- Hop routing (OLSR).
  - ✓ Rate guarantee (PM) or End to End Guaranteed Delivery (commands, etc.)

## TFC Scheduling by Service type

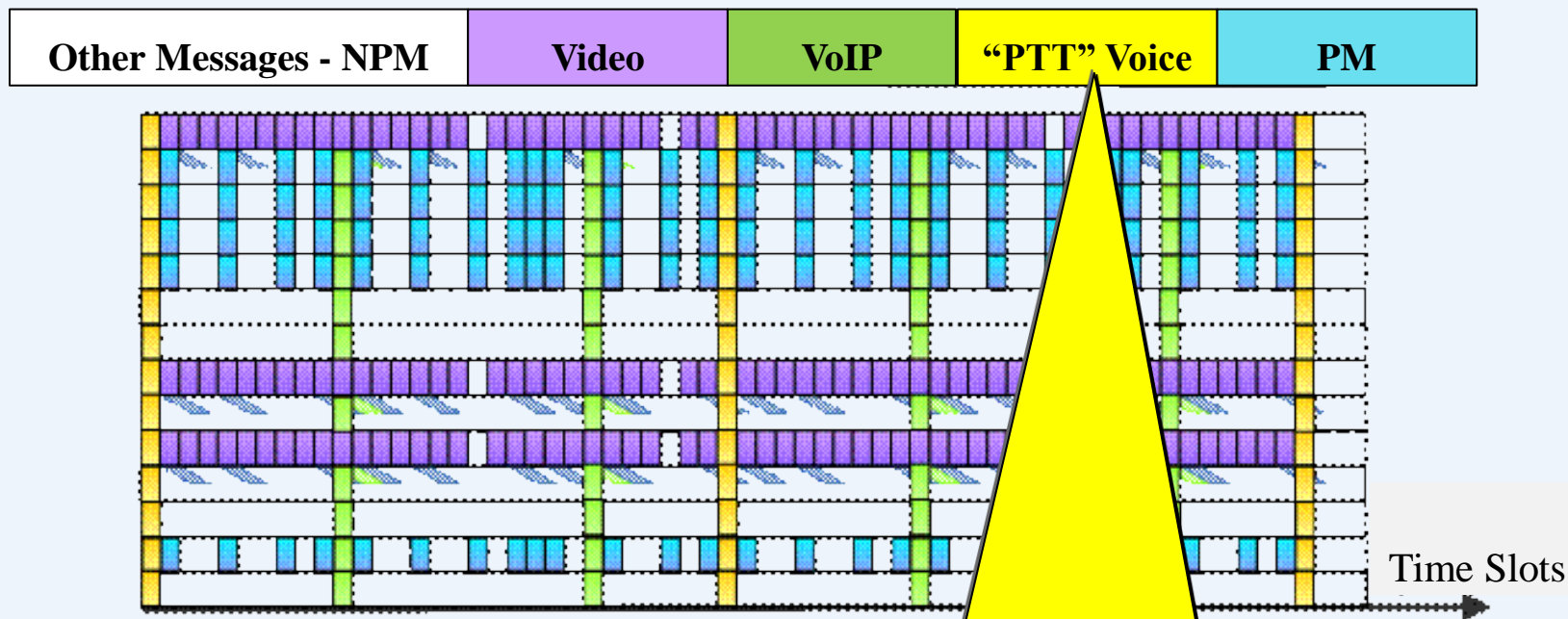


### Periodical Refreshing Messages - for SA, CA, and control messages (Area context)

- Delay sensitive, need Assured message rate
- Adjustable reliability using FEC.
- Broadcast N- Hop Relaying.



## TFC Scheduling by Service type

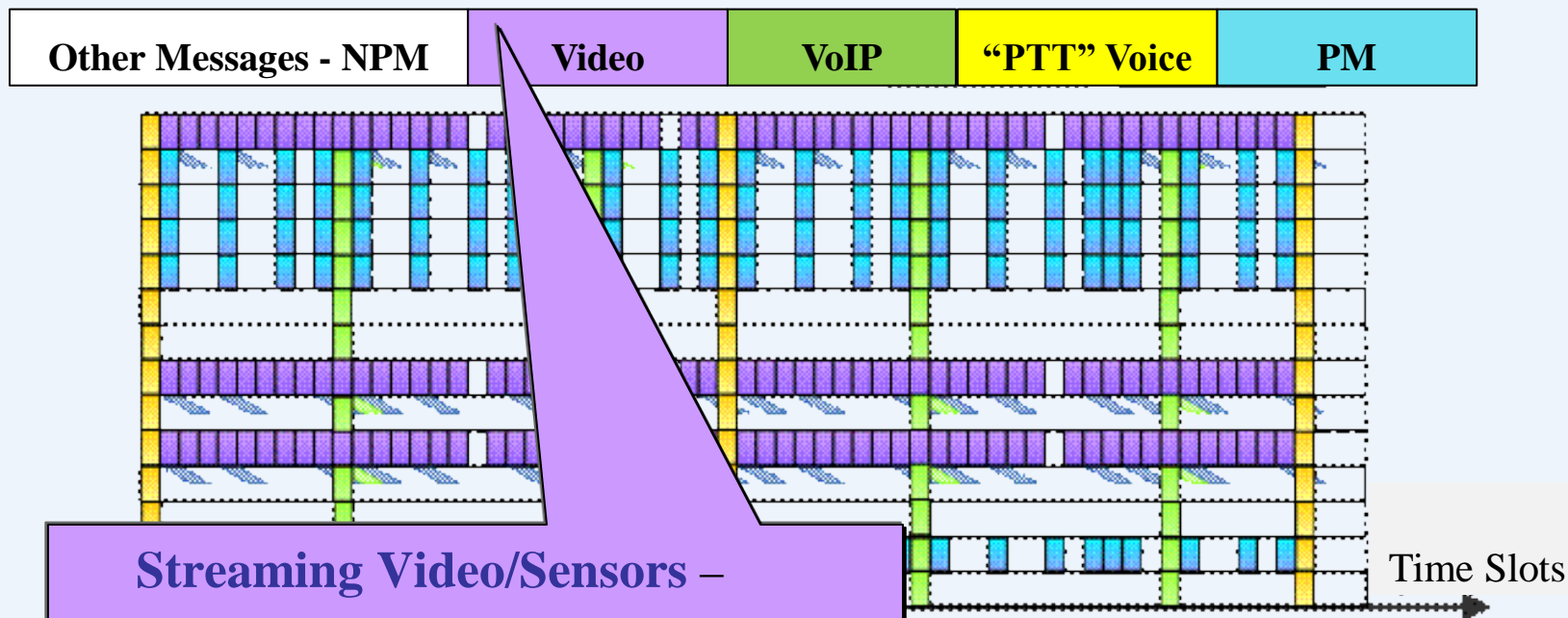


### “PTT Combat Voice” - Critical Delay.

- **Dedicated channels and time-slots.**
- 1 Channel = Several Voice Multicast Groups. (Since 1 voice session is ~ 4 Kb/sec, a 1 Mb/sec channel can supports several sessions).
- Broadcast 1-Hop Relaying (OLSR).



## TFC Scheduling by Service type

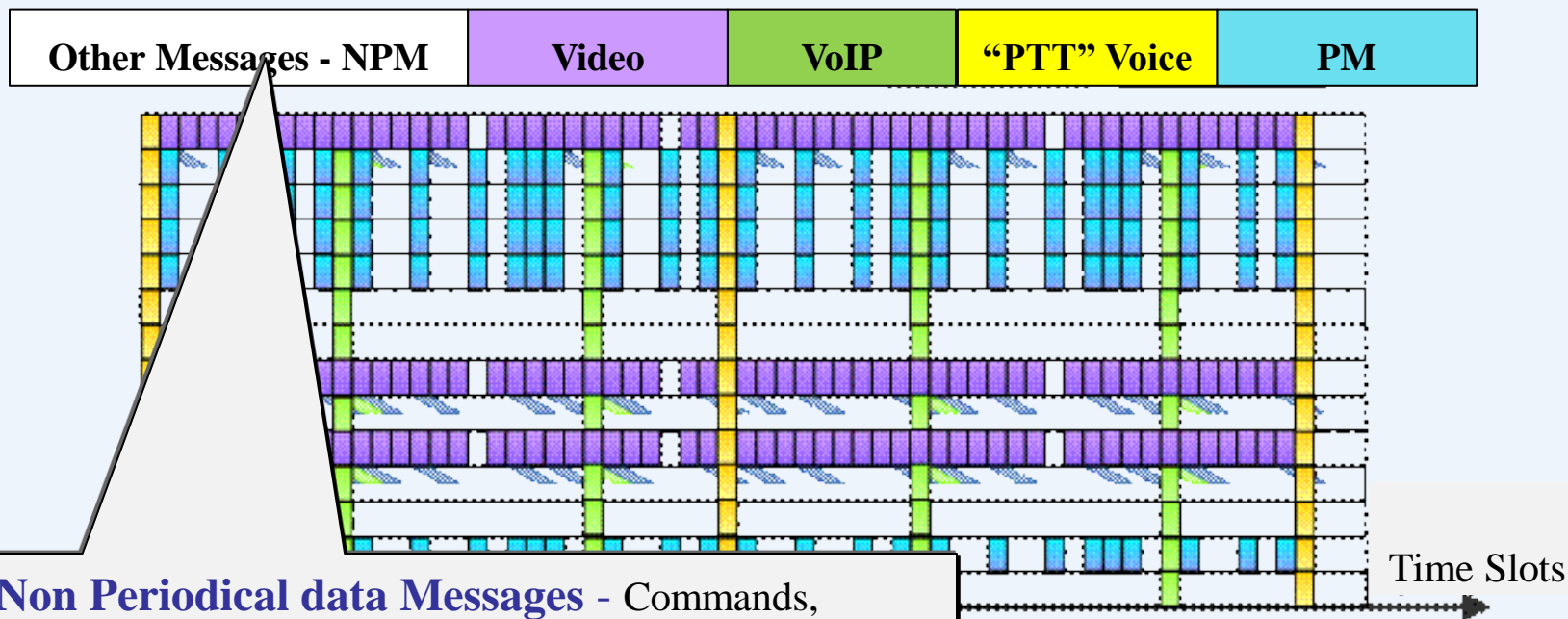


### Streaming Video/Sensors –

Relatively few Sensor nodes per area.

- Can suffer small delay.
- Depending on compression and resolution, 0.25 – 1 Mb/s per stream => 1 MHz channel per sensor.
- UC, MC or BC. Up to 1 Hop relaying.
- Time Slots “Punctured” for Voice + PM reception (at reduced rate).

## TFC Scheduling by Service type

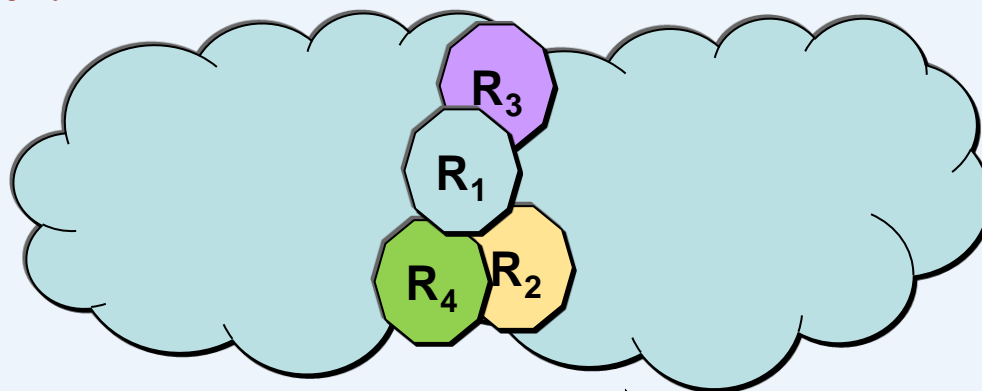


**Non Periodical data Messages** - Commands, Reports etc. (**Units or Functional group context**)

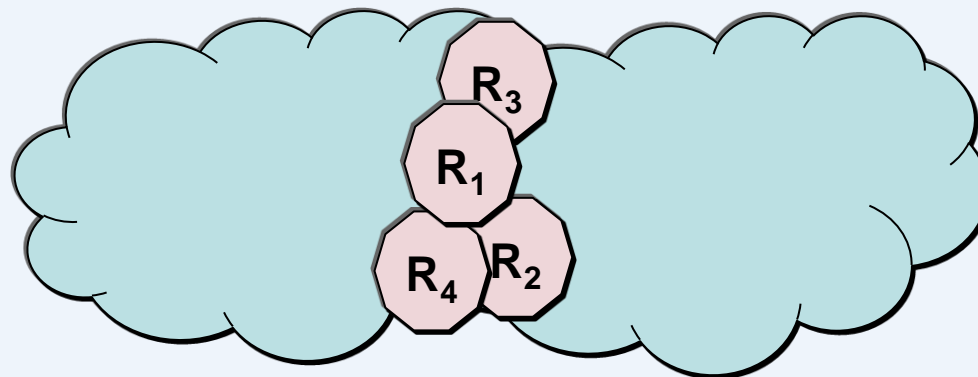
- On-demand variable-length messages, Unicast or Multicast.
- Delay insensitive.
- Many-Hops routing + Delivery Guarantee is needed.
- TFCs allocation is in all TFCs not constrained by PM, Voice, Video.

## Routing advantages of MCR

- **Single Channel Manet** –  $R_1$ , the only “blue channel” node that Rx both sides of its net, is a bottleneck Relay; although  $R_2 \div R_4$  are within 1-HOP range of both sides of the blue net, they aren’t tuned it and can’t relay. **Relay Bottleneck!**

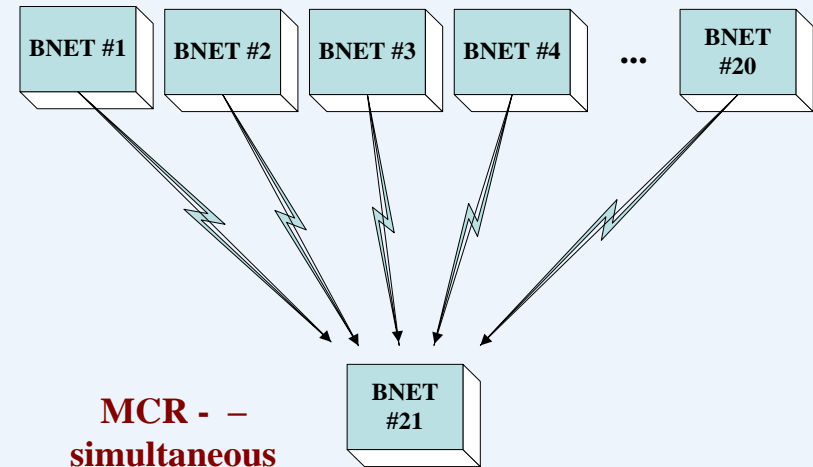


- **MCR Manet** – All nodes Rx all channels  $\Rightarrow$  4 nodes can relay – **No bottleneck, better flow-control**



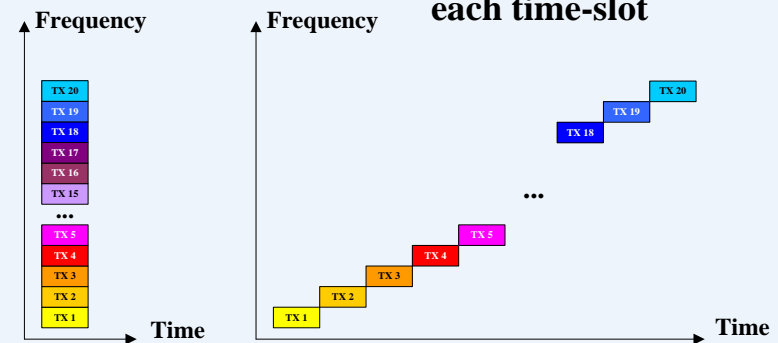
# Summary - Advantages of Multi Channel Reception MANET

- A breakthrough in network capacity, scalability and delays
- Less overhead - less Relays, Negotiations, Reservations
- Simpler network planning
- Quick force merge, split & pass-through – Reconfiguration not needed.
- More performance in much less hardware
- Cognitive ready – each SDR is a “continuous” Spectrum Analyzer



**MCR - –  
simultaneous  
reception of >  
100 independent  
channels**

**Other modern MANET –  
one received channel at  
each time-slot**



# The BNET Multi-Channel Reception SDR Family

*Breakthrough network Capacity, Delay, Scalability - Data, Video, Voice on the move*



## BNET- Airborne

4 Rx Bands in 30-2000  
 2X50 W Internal Transmitters  
 Manet +AM/FM+ UHF Satcom  
 11 Kg, ½ ATR (35X19X13 cm)  
 Tx =Rate: up to 10 Mb/s  
 Rx Rate: > 200 Mb/s !

## BNET- Vehicular

2 selectable Rx Bands  
 1 50 W Tx  
 Manet + AM/FM  
 8 Kg, 30X20X10 cm  
 Tx =Rate: up to 10 Mb/s  
 Rx Rate: > 200 Mb/s

## BNET- Hand Held

1 selectable band  
 Manet or AM/FM  
 5 W Tx, 1.2 Kg  
 Tx =Rate: up to 2 Mb/s  
 Rx Rate: 20 Mb/s  
 (A 4 Kg, 2 Bands, 20 W, larger battery Manpack will also be available)

- Different Chasses & Cooling
- Identical RF, DPU, SBC Modules



# Thanks For Your Attention !

## Questions?

Please call {mikiw, michalpr}@rafael.co.il





## References

- [1] J. H. Ju and V. O. K. Li, “TDMA Scheduling Design of Multihop Packet Radio Networks Based on Latin Squares”, In IEEE Journal on Selected Areas in Communications, Volume 17, Issue 8, Page(s):1345-1352, Aug. 1999
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- [3] L. Bao and J.J. Garcia-Luna-Aceves. “A New Approach to Channel Access Scheduling for Ad Hoc Networks”. In Proc. ACM Seventh Annual International Conference on Mobile Computing and networking, Rome, Italy, Jul. 2001
- [4] R. Rozovsky and P. R. Kumar. “SEEDEX: a MAC protocol for ad hoc networks”. In Proc. of the 2nd ACM International Symposium on Mobile Ad Hoc Networking and Computing (MobiHoc) 2001, pages 67–75, Long Beach, CA, USA, Oct. 4-5 2001.
- [5] M. Wermuth, Y. Wermuth, M. Weiss, S. Avadis, Y. Fuchs, “An Improved Mobile Ad-Hoc Network”, Pending Patent application # IL216282, Dec 2011.
- [6] S. Avadis, A. Lerner, Z. Nutov, “MMM: multi-channel TDMA with MPR capabilities for MANETs”, Wireless networks magazine, Vol. 18 No. 5, July 2012. Available online from Springer as DOI 10.1007/s11276-012-0468-6.

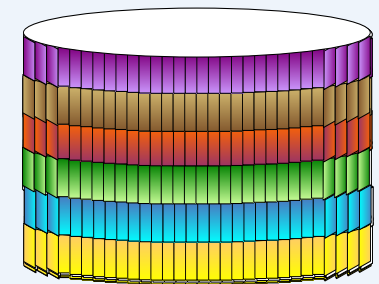
# Backup Slides

# The BNet MANET WF

## Mitigating the Spectrum challenge

- **Easy and simple pre-planning:**

- Nets are Multicast groups, not frequencies
- Given a spectrum allocation and QoS requirements, The net dynamically utilizes only the needed channels to meet the QoS requirements efficiently



- **Highly efficient automatic multi-rate capabilities**

- Utilize the ultra wideband receiver as a rate sensor
- Automatically adapt the data rate (from BPSK/QPSK to 64QAM )